A Transformational Approach to Scientific Software: The Mathematics of Array (MoA) Fast Fourier Transform (FFT) with OpenACC

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Abstract

We extend a methodology for designing efficient parallel and distributed scientific software for GPUs introduced in Rosenkrantz *et al.* [1]. This methodology utilizes sequences of mechanizable algebra-based optimizing transformations. Starting from a high-level algebraic algorithm description in "*A Mathematics of Arrays*" (MoA) [2], abstract multiprocessor plans are developed and refined to specify which computations are to be done by each processor. Starting with the OpenMP program in Fortran 90 produced in Rosenkrantz *et al.* [1], we extend it to include OpenACC for GPU support. Our studies show what is needed in OpenACC, support in Fortran compilers for GPUs, and what issues we encountered and resolved.

[1] Harry B. Hunt, Lenore R. Mullin, Daniel J. Rosenkrantz, and James E. Raynolds (2008), "A Transformation-Based Approach for the Design of Parallel and Distributed Scientific Software: The FFT" <u>arXiv:0811.2535</u>

[2] Lenore R. Mullin (1988), "A Mathematics of Arrays", PhD dissertation. Syracuse University. doi:10.5555/915213

What is Mathematics of Arrays (MoA)?

- Mathematics of Arrays (MoA) is a theory of arrays, where everything can be written in terms of <u>array shapes</u> and *index composition* ("<u>Psi</u>" <u>reduction</u>).
- MoA makes an *ideal* formulation of computation when combined with Lambda calculus [3].
- Both MoA and Lambda calculus possess the **Church-Rosser property** [4].
- [3] Klaus Berkling, "Arrays and the Lambda Calculus", (1990), Technical Report 93, Syracuse University. https://surface.syr.edu/eecs_techreports/93/

[4] Benjamin Chetioui, Lenore Mullin, Ole Abusdal, Magne Haveraaen, Jaakko Järvi, and Sandra Macià (2019), "Finite difference methods fengshui: alignment through a mathematics of arrays", ARRAY 2019: Proceedings of the 6th ACM SIGPLAN International Workshop on Libraries, Languages and Compilers for Array Programming. doi:10.1145/3315454.3329954

What is Mathematics of Arrays (MoA)? (cont'd)

- MoA syntax is heavily inspired by the array algebra in **Ken Iverson**'s **APL programming language** [5].
- When designing and implementing the APL machine,
 Phil Abrams introduced the idea of using array shapes [6].
- Lenore Mullin added *closure* on the algebra [2].

[2] Lenore R. Mullin (1988), "A mathematics of arrays", PhD dissertation. Syracuse University. doi:10.5555/915213

[5] Kenneth E. Iverson (1980), "Notation as a tool of thought", 1979 ACM Turing Award lecture, Communications of the ACM, 23 (8), 444–465. doi:10.1145/358896.358899

[6] Philips S. Abrams, "An APL machine" (1970), technical report SLAC-114. Stanford University. https://www.slac.stanford.edu/pubs/slacreports/reports07/slac-r-114.pdf

The Cooley–Tukey Fast Fourier Transform (FFT) Algorithm

From van Loan [7], p. 46, Algorithm (1.6.2):

$$\begin{aligned} x \leftarrow P_n x & B_L = \begin{bmatrix} I_{L_0} & \Omega_{L_0} \\ I_{L_0} & -\Omega_{L_0} \end{bmatrix} & \\ L \leftarrow 2^t; r \leftarrow n/L & L = 2^q & \Omega_{L_0} = \begin{bmatrix} 1 & \\ \omega_L & \\ & \ddots & \\ x_{L \times r} \leftarrow B_L x_{L \times r} & r = n/L \\ & L_0 = L/2 & & \end{bmatrix} \\ end & end & L = L/2 & L = L/2 & L = L/2 \\ \end{aligned}$$

[7] Charles van Loan (1992), "Computational Frameworks for the Fast Fourier Transform", Frontiers in Applied Mathematics, SIAM. <u>doi:10.1137/1.9781611970999</u>

The Cooley–Tukey Fast Fourier Transform (FFT) Algorithm (cont'd)

In principle, any FFT algorithm could be chosen, but we started with this particular one because:

- It uses **arrays** effectively.
- It uses the **Kronecker product** (= outer product in MoA).
- It is a **simple** operation (1-dimensional radix-2 FFT) that can be easily extended to *N* dimensions.
- It is one of the **most popular** algorithms for the FFT.
- Van Loan is a respected expert on arrays.

MoA and Transformations

- MoA uses a systematic, algebraic design methodology to reduce the FFT algorithm into a semantic <u>denotational normal form</u> (DNF).
- Next, we employ "*dimension lifting*" to get the operational normal form (ONF).
 - It describes the problem and *partitioning* of *data* over *processors* and *memory hierarchies*.
 - It allows one to mathematically prove the **efficiency** and **correctness** of a given algorithm as measured in terms of a set of metrics (such as processor/network/memory speeds).
- Such an approach allows the average programmer to achieve *high-level optimizations* similar to those used by compiler writers, e.g. the notion of tiling.
- We envision scientific programs in the future will be developed in an *interactive* development environment that combines *human judgment* with *compiler-like analysis*, such that transformation and verification of its correctness can be done *mechanically*.

The bit reversal permutation

- Notice the first step of the Cooley-Tukey algorithm is the **bit reversal** permutation $x \leftarrow P_n x$.
- van Loan [7], p. 39, Algorithm (1.5.1): For $t \ge 1$ and $0 \le k < 2^t$,

$$f \leftarrow 0; m \leftarrow k$$

for $q = 0: t - 1$
 $s \leftarrow \text{floor}(m/2)$

$$\{b_q = m - 2s\}$$
$$j \leftarrow 2j + (m - 2s)$$

 $m \gets s$

- We developed this algorithm with MoA: $P_n = \vec{i} \psi \bigotimes \left(\left(\langle t \rangle \rho 2 \right) \rho \left(\iota 2^t \right) \right) \qquad 0 \stackrel{*}{\leq} \vec{i} \stackrel{*}{\leq} \left(\langle t \rangle \rho 2 \right)$
- The implementation uses bit shifts:

```
revivec = 0; k = i
D0 j = 1, t
revivec = ISHFT( revivec, 1 )
revivec = IOR( revivec, IAND( k, 1_dl ) )
k = ISHFT( k, -1 )
END D0
revivec = k
```

end

[7] Charles van Loan (1992), "Computational Frameworks for the Fast Fourier Transform", Frontiers in Applied Mathematics, SIAM. <u>doi:10.1137/1.9781611970999</u>

The base CPU-only version

- Rosenkrantz *et al.* preprint [1] details the rationale and development from the original Cooley-Tukey algorithm to the MoA version.
- The Operational Normal Form (ONF) from MoA enables one to choose the block size that gives the **best performance** for any individual machine -assuming intentional information can be processed by a compiler.
- The generic design of the program is then implemented in Fortran with three versions: **sequential**, shared memory (**OpenMP**), and distributed memory (**MPI**).

 [1] Harry B. Hunt, Lenore R. Mullin, Daniel J. Rosenkrantz, and James E. Raynolds (2008),
 "A Transformation-Based Approach for the Design of Parallel and Distributed Scientific Software: The FFT", <u>arXiv:0811.2535</u>

The GPU-enabled version with OpenACC

- The GPU has its *unique* memory hierarchy, which is arguably a *mix* between shared memory and distributed memory.
 - Device memory is distinct from host memory (unless when using unified memory, which we chose <u>not</u> to use).
 - Global device memory is shared between execution units (SMs on NVIDIA, CUs on AMD, etc.)
 - On NVIDIA GPU, CUDA shared memory is shared between threads in the same block.
- For this reason, we start with the **shared memory model** and the OpenMP version of the code, and port it for GPUs using OpenACC.

Why use OpenACC?

- **Ease of use**: just "sprinkle" some directives at the right places, and the code is GPU-enabled!
- Ability to maintain a **single source code** for both the CPU-only and the GPU-enabled versions.
- Portability: same code can run on both NVIDIA GPU (pgf90/nvfortran) and AMD GPU (gfortran ≥ 10). Rumor has it that Intel GPU support is coming soon too.

Favorite OpenACC features

- Full Fortran support (unlike HIP or DPC++)
- Manual control over grid and block size via

num_gangs, num_workers, and vector_length clauses

!\$acc parallel num_gangs(ngang) num_workers(nworker) vector_length(veclen)

116, Generating Tesla code
 121, !\$acc loop gang(1024) ! blockidx%x
 123, !\$acc loop worker(8) ! threadidx%y
 125, !\$acc loop vector(128) ! threadidx%x

Favorite OpenACC features (cont'd)

 Competitive *performance* w.r.t. CUDA on NVIDIA GPUs

Read: log2r :	ing times [.] n init	-moafft-si: h2d	ze-210907. bitrev	dat : fft_fwd	d2h
20 (21 (0.007540 0.015312	0.030121 0.033479	0.000293 0.000464	0.009980 0.011414	0.003912
22 (0.029750	0.039017	0.000761	0.012614	0.015826
23 (24 (0.060/49	0.050531	0.001401	0.014182	0.032252
:					
Reading times-cufft-size-210908.dat :					
log2r :	n init	h2d	plan	fft_fwd	d2h
20 (0.009177	0.003474	0.001058	0.000209	0.003718
21 (0.018347	0.006826	0.001065	0.000333	0.007214
22 (0.035533	0.013427	0.001902	0.000578	0.014378
23 0	0.071028	0.027496	0.002069	0.001195	0.028409
24 (:	0.147658	0.054567	0.002066	0.002322	0.057793



Favorite OpenACC features (cont'd)

cache directive for using CUDA shared memory

```
!$acc loop vector collapse(2) private( c, d ) independent
do i = 0, size-1, Lnew
do j = 0, (Lnew/2)-1
   !$acc cache( weight_p, zblock_p )
   c = weight_p(j) * zblock_p( i + j + Lnew/2 )
   d = zblock_p( i + j )
   zblock_p( i + j ) = d + c
   zblock_p( i + j + Lnew/2 ) = d - c
end do ! j
end do ! i
```

• Array reduction over an arbitrary axis (supposedly introduced in OpenACC 2.7, but we can't seem to get it to work...)

```
integer :: i, j, m, n
real, allocatable :: vec(:), mat(:,:)
allocate( vec(m), mat(m,n) )
!$acc data create( vec, mat )
!$acc parallel
!$acc loop gang
do i = 1, m
 vec(i) = 0.0
 !$acc loop vector reduction(+:vec)
 do j = 1, n
   vec(i) = vec(i) + mat(i,j)
 end do
end do
```

- Array reduction over an arbitrary axis (supposedly introduced in OpenACC 2.7, but we can't seem to get it to work...)
- Manual *synchronization* routines

```
__syncwarp()
__syncthreads()
cudaStreamSynchronize() == !$acc wait
cudaDeviceSynchronize()
```

- Array reduction over an arbitrary axis (supposedly introduced in OpenACC 2.7, but we can't seem to get it to work...)
- Manual *synchronization* routines
- No "hidden" cost of OpenACC runtime library initialization

```
! Notice we have moved OpenACC runtime
! initialization (~300 ms) outside of
! the timing region for initialization
!Sacc init
t0 = now()
do i = 0, n - 1
  z(i) = cmplx( i, 0._dd, kind=dz )
end do
t1 = now()
time init = t1 - t0
```

- Array reduction over an arbitrary axis (supposedly introduced in OpenACC 2.7, but we can't seem to get it to work...)
- Manual *synchronization* routines
- No "hidden" cost of OpenACC runtime library initialization
- Support for *pinned memory*

```
! CUDA Fortran
ATTRIBUTES(pinned) :: x
```

OpenMP target offload?

- Once compiler support matures for OpenMP target offload, we will consider porting from OpenACC into OpenMP target offload.
- Doing this would enable *concurrent parallelism* for both multicore CPUs and accelerators (GPUs) under the *same programming model*.
- Automated *runtime translation* tools (e.g. clacc and flacc [8]) are welcome too!

[8] <u>https://csmd.ornl.gov/project/clacc</u>

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